

# In-Degree and PageRank of Web pages: Why do they follow similar power laws?

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## Abstract

PageRank is a popularity measure designed by Google to rank Web pages. Experiments confirm that PageRank values obey a power law with the same exponent as In-Degree values. This paper presents a novel mathematical model that explains this phenomenon. The relation between PageRank and In-Degree is modelled through a stochastic equation, which is inspired by the original definition of PageRank, and is analogous to the well-known distributional identity for the busy period in the  $M/G/1$  queue. Further, we employ the theory of regular variation and Tauberian theorems to analytically prove that the tail distributions of PageRank and In-Degree differ only by a multiple factor, for which we derive a closed-form expression. Our analytical results are in good agreement with experimental data.

## Categories and Subject Descriptors

H.3.3:[**Information Storage and Retrieval**]: Information Search and Retrieval–*Retrieval models*; G.3:[**Mathematics of Computing**]: Probability and statistics – *Stochastic processes, Distribution functions*

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# 1 Introduction

In this paper we study the relation between the probability distributions of the PageRank and the In-Degree of a randomly selected Web page. The notion of *PageRank* was introduced by Google in order to numerically characterize the popularity of Web pages. The original description of PageRank presented in [11] is as follows:

$$PR(i) = c \sum_{j \rightarrow i} \frac{1}{d_j} PR(j) + (1 - c), \quad (1)$$

where  $PR(i)$  is the PageRank of page  $i$ ,  $d_j$  is the number of outgoing links of page  $j$ , the sum is taken over all pages  $j$  that link to page  $i$ , and  $c$  is the “damping factor”, which is some constant between 0 and 1. The *In-Degree* of a Web page denotes simply the number of incoming hyperlinks to that page. From equation (1) it is clear that the PageRank of a page depends on its In-Degree and the importance (i.e. PageRanks) of the pages that link to it.

We focus in particular on the *tail asymptotics* for PageRank and its connection to In-Degree. By tail of the PageRank distribution we mean the fraction of pages  $\mathbb{P}(PR > x)$  having PageRank greater than  $x$ , where  $x$  is large. Thus, we are concerned only with pages of high ranking. A common way to analyze tail behavior is to find an asymptotic expression  $p(x)$  such that  $\mathbb{P}(PR > x)/p(x) \rightarrow 1$  as  $x \rightarrow \infty$ . In this case,  $p(x)$  and  $\mathbb{P}(PR > x)$  are asymptotically similar, and thus, we can approximate  $\mathbb{P}(PR > x)$  by  $p(x)$  for large enough  $x$ .

Pandurangan et al. [22] observed that the tails of PageRank and In-Degree distributions for Web data seem to follow power laws with the same exponent. Loosely speaking, a power law with exponent  $\alpha$  means that the probability that the random variable takes values greater than some large number  $x$  is approximately proportional to  $x^{-\alpha}$ . Formally, this can be modelled as asymptotic similarity of PageRank and In-Degree tail distributions to some power law functions. It turns out that for both PageRank and In-Degree distributions, the power law exponent  $\alpha$  is about 1.1 for cumulative plots, which gives the famous value 2.1 for the density.

Recent extensive experiments by Donato et al. [12] and Fortunato et al. [16] confirmed the similarity in tail behavior observed in [22]. Becchetti and Castillo [6] extensively investigated the influence of the damping factor  $c$  on the power law behavior of the PageRank. They have shown that the PageRank of the top 10% of the nodes always follows a power law with the same exponent independent of the value of the damping factor. Our own experiments based on Web data from [1] are also in agreement with [22] (see Figure 1 in Section 5.2 and the discussion there).

Obviously, equation (1) suggests that PageRank and In-Degree are intimately related, but this formula by itself does not explain the observed similarity in tail behavior. Furthermore, the linear algebra methods that have been commonly used in the PageRank literature [7, 19] and proved very successful for designing efficient computational methods, seem to be insufficient for modelling and analyzing the asymptotic properties of the PageRank distribution.

The goal of our paper is to provide mathematical evidence for the power-law behavior of PageRank and its relation to the In-Degree distribution. We propose a stochastic model that aims to explain this relation. Our approach is inspired by the techniques from applied probability and stochastic operations research. The relation between PageRank and In-Degree is modelled through a distributional identity which is analogous to the equation for the busy period in the M/G/1 queue (see e.g. [23]). Further, we analyze our model using the approach employed in [20] for studying the tail behavior of the busy period in case the service times are regularly varying random variables. This fits in our research because regular variation is in fact a generalization of the power law, and it has been widely used in queueing theory to model self-similarity, long-range dependence and heavy tails [24]. Thus, we use the notion of regular variation to model the power law distribution of the In-Degree. For the sake of completeness, in Section 2, we will introduce regularly varying random variables and describe their basic properties.

To obtain the tail behavior of PageRank in our model, we use Laplace-Stieltjes transforms and apply Tauberian theorems presented in the paper by Bingham and Doney [8], see also Theorem 8.1.6 in [9]. Even though our model contains some rather rigid simplifying assumptions – the most notable being independence between pages that link to the same page and a constant Out-Degree for all pages – these techniques allow us to prove the similarity in tail behavior for PageRank and In-Degree, thus suggesting that our assumptions do not touch upon the underlying reasons for this similarity. Moreover, our analysis allows to explicitly derive the constant multiple factor that quantifies the difference between PageRank and In-Degree tail behavior. We describe the model in Section 3 and provide the main result and its proof in Section 4. The technical proofs of ancillary statements are deferred to the Appendix. As discussed in Section 5, our analytical results show a good agreement with real Web data.

We believe that our approach is extremely promising for analyzing the PageRank distribution and solving other problems related to the structural properties of the Web. At the end of this paper, we will briefly mention other possibilities for probabilistic analysis of the PageRank distribution. In particular, in Section 5.3, we provide experimental results for Growing Networks [2], and in Section 6, we draw a parallel between the recent studies [4, 15] on PageRank behavior in this class of graph models and our present work.

## 2 Preliminaries

This section describes important properties of regularly varying random variables. We follow definitions and notations by Bingham and Doney [8], Meyer and Teugels [20] and Zwart [24]. More comprehensive details can be found in [9].

**Definition 1.** A function  $V(x)$  is regularly varying of index  $\alpha \in \mathbb{R}$  if for

every  $t > 0$ ,

$$\frac{V(tx)}{V(x)} \rightarrow t^\alpha \quad \text{as } x \rightarrow \infty.$$

If  $\alpha = 0$ , then  $V$  is called *slowly varying*.

**Definition 2.** Function  $L$  is slowly varying if for every  $t > 0$ ,

$$\frac{L(tx)}{L(x)} \rightarrow 1 \quad \text{as } x \rightarrow \infty.$$

A function  $V(x)$  is regularly varying if and only if it can be written in the form

$$V(x) = x^\alpha L(x),$$

for some slowly varying  $L(x)$ .

The following lemma provides a useful bound for slowly varying functions.

**Lemma 1.** (*Potter bounds*) Let  $L$  be a slowly varying function. Then, for any fixed  $A > 1, \delta > 0$  there exists a finite constant  $K > 1$  such that for all  $x_1, x_2 > K$ ,

$$\frac{L(x_1)}{L(x_2)} \leq A \max \left\{ \left( \frac{x_1}{x_2} \right)^\delta, \left( \frac{x_1}{x_2} \right)^{-\delta} \right\}.$$

**Definition 3.** In probability theory, a random variable  $X$  is said to be regularly varying with index (or exponent)  $\alpha$  if its distribution  $F$  is such that

$$1 - F(x) \sim x^{-\alpha} L(x) \quad \text{as } x \rightarrow \infty,$$

for some positive slowly varying function  $L(x)$ . Here, as in the remainder of this paper, the notation  $a(x) \sim b(x)$  means that  $a(x)/b(x) \rightarrow 1$ .

Denote by  $f(s) = \mathbb{E}e^{-sX}$ ,  $s > 0$ , the Laplace-Stieltjes transform of  $X$ , and let  $\xi_n = \int_0^\infty x^n dF(x)$  be the  $n$ th moment of  $X$ . The successive moments of  $F$  can be obtained by expanding  $f$  in a series at  $s = 0$ . More precisely, we have the following.

**Lemma 2.** The  $n$ th moment of  $X$  is finite if and only if there exist numbers  $\xi_0 = 1$  and  $\xi_1, \dots, \xi_n$ , such that

$$f(s) - \sum_{i=0}^n \frac{\xi_i}{i!} (-s)^i = o(s^n) \quad \text{as } s \rightarrow 0.$$

If  $\xi_n < \infty$  then we introduce the notation ( $n \in \mathbb{N}$ )

$$f_n(s) = (-1)^{n+1} \left( f(s) - \sum_{i=0}^n \frac{\xi_i}{i!} (-s)^i \right). \quad (2)$$

**Remark 1.** *It follows from Lemma 2 that the  $n$ th moment of  $X$  is finite if and only if there exist numbers  $\xi_0 = 1$  and  $\xi_1, \dots, \xi_n$  such that  $f_n(s) = o(s^n)$  as  $s \rightarrow 0$ .*

The following theorem establishes the relation between asymptotic behavior of regularly varying distribution and its Laplace-Stieltjes transform. This result plays an essential role in our analysis.

**Theorem 1.** *(Tauberian Theorem) If  $n \in \mathbb{N}$ ,  $\xi_n < \infty$ ,  $\alpha = n + \beta$ ,  $\beta \in (0, 1)$ , then the following are equivalent*

$$(i) \ f_n(s) \sim (-1)^\alpha \Gamma(1 - \alpha) s^\alpha L\left(\frac{1}{s}\right) \quad \text{as } s \rightarrow 0,$$

$$(ii) \ 1 - F(x) \sim x^{-\alpha} L(x) \quad \text{as } x \rightarrow \infty.$$

Here and in the remainder of the paper we use the letter  $\alpha$  to denote the index of a complementary distribution function  $1 - F(x)$  rather than a density. The power law exponent of the In-Degree in the Web graph then becomes 1.1 rather than 2.1.

### 3 The model

In this section we introduce a model that describes the relation between PageRank and In-Degree in the form of a stochastic equation. This model naturally follows from the definition of PageRank in (1), and is analytically tractable, thus enabling us to obtain the asymptotic behavior of PageRank. As will become clear, we make several rather severe simplifying assumptions. Nevertheless, the theoretical results of this model show a good match with observed Web graph behavior.

#### 3.1 Relation between In-Degree and PageRank

Our goal now is to describe the relation between PageRank and In-Degree. To this end, we keep equation (1) almost unchanged but we make several assumptions. First, let  $R$  be the PageRank of a randomly chosen page. We treat  $R$  simply as a random variable whose distribution we want to determine. Second, we assume that the number of outgoing links  $d \geq 1$  is the same for each page, and we do not consider the influence of pages without outgoing links ('dangling nodes'). Although these assumptions are not realistic, they help us to focus on the influence of In-Degree, without considering other factors. We note however that present model allows for various generalizations. For instance, we can account for the dangling nodes as discussed at the end of Section 4.

Under the assumptions above, the random variable  $R$  satisfies a distributional identity

$$R \stackrel{d}{=} c \sum_{j=1}^M \frac{1}{d} R_j + (1 - c), \quad (3)$$

where  $M$  is the In-Degree of the considered random page.

We now make an assumption that  $M$  and the  $R_j$ 's are independent, and  $R_j$ 's have the same distribution as  $R$  itself. We note that the independence assumption is obviously not true in general. However, it is also not the case that the PageRank values of the pages linking to the same page  $i$  are directly related, so we may assume independence in this study.

The novelty of our approach is that we treat the PageRank as a random variable which solves a certain stochastic equation. However, this approach is quite natural if our goal is to explain the power law behavior of PageRank because the 'power law' is merely a description of a certain class of probability distributions.

One of the nice features of stochastic equation (3) is that it has the same form as the original formula (1). Thus, we may hope that our model correctly describes the relation between In-Degree and PageRank. This is easy to verify in the extreme (unrealistic) case when all pages have the same In-Degree  $d$ . In this situation, the PageRanks of all pages are equal, and it is easy to see that  $R \equiv 1$  constitutes the unique solution of (3).

### 3.2 In-Degree Distribution

It is well-known that the In-Degree of Web pages follows a power law. For our analysis however we need a more formal description of this random variable, thus, we suggest to employ the theory of regular variation. We model the In-Degree of a randomly chosen page as a nonnegative, integer, regularly varying random variable, which is distributed as  $N(X)$ , where  $X$  is regularly varying with index  $\alpha$  and  $N(x)$  is the number of Poisson arrivals on the time interval  $[0, x]$ . Without loss of generality, we assume that the rate of the Poisson process is equal to 1.

The advantage of this construction is that we do not need to impose any restrictions on  $X$  and at the same time ensure that the In-Degree is integer. We claim that the random variable  $N(X)$  will also be regularly varying with the same index as  $X$ , or, more informally,  $N(X)$  follows a power law with the same exponent. Thus, we can think of  $N(X)$  as the In-Degree of a random Web page. For the sake of completeness we present the formal statement and its proof in the remainder of this section.

Let  $F_X$  and  $F_{N(X)}$ ,  $f$  and  $\phi$  be the distribution functions and the Laplace-Stieltjes transforms of  $X$  and  $N(X)$ , respectively. Since the random variable  $X$  is regularly varying, we have by definition

$$1 - F_X(x) \sim x^{-\alpha} L(x) \text{ as } x \rightarrow \infty, \quad (4)$$

where  $L(x)$  is some slowly varying function. Then we will claim that for  $N(X)$  the following also holds:

$$1 - F_{N(X)}(x) \sim x^{-\alpha} L(x) \text{ as } x \rightarrow \infty. \quad (5)$$

For completeness, we prove this statement using the Tauberian theorem (Theorem 1). To this end, we first have to show that the corresponding moments

of  $X$  and  $N(X)$  always exist together. Assuming that  $\mathbb{E}X = d$  we immediately get  $\mathbb{E}N(X) = d$ . Next, consider the generating function of  $N(X)$ ,

$$\mathbb{G}_{N(X)}(s) := \mathbb{E}s^{N(X)} = \int_0^\infty \mathbb{E}s^{N(t)} dF_X(t) = \int_0^\infty e^{-t(1-s)} dF_X(t) = f(1-s), \quad (6)$$

from which we derive the Laplace-Stieltjes transform of  $N(X)$  in terms of the Laplace-Stieltjes transform of  $X$ :

$$\phi(w) = \mathbb{E}e^{-wN(X)} = f(1 - e^{-w}).$$

Now, denote by  $\xi_1 = d, \xi_2, \dots, \xi_n$  and  $\nu_1 = d, \nu_2, \dots, \nu_n$  the first  $n$  moments of  $X$  and  $N(X)$ , respectively, and define  $\xi_0 = \nu_0 = 1$ . Here  $n$  is the largest integer smaller than  $\alpha$ , and thus  $\xi_n$  is the highest finite moment of  $X$ . Then we can establish an auxiliary result formulated in the next lemma (see the Appendix for the proof).

**Lemma 3.** *The following are equivalent*

- (i)  $\xi_n < \infty$ ,
- (ii)  $\nu_n < \infty$ .

**Remark 2.** *If we define*

$$f_n(s) = (-1)^{n+1} \left( f(s) - \sum_{i=0}^n \frac{\xi_i}{i!} (-s)^i \right) \text{ and}$$

$$\phi_n(s) = (-1)^{n+1} \left( \phi(s) - \sum_{i=0}^n \frac{\nu_i}{i!} (-s)^i \right)$$

as in (2), then we can reformulate Lemma 3 as follows:

$$f_n(s) = o(s^n) \quad \text{if and only if} \quad \phi_n(s) = o(s^n).$$

Now, we can use Theorem 1 to prove that (4) implies (5). In fact the reverse also holds, as stated in the following theorem.

**Theorem 2.** *The following are equivalent*

- (i)  $1 - F_X(x) \sim x^{-\alpha} L(x)$  as  $x \rightarrow \infty$ ,
- (ii)  $1 - F_{N(X)}(x) \sim x^{-\alpha} L(x)$  as  $x \rightarrow \infty$ .

*Proof.*

(i)  $\rightarrow$  (ii) From Theorem 1 for  $X$  we know that

$$1 - F_X(x) \sim x^{-\alpha} L(x), \quad x \rightarrow \infty$$

implies

$$f_n(t) \sim (-1)^\alpha \Gamma(1 - \alpha) t^\alpha L\left(\frac{1}{t}\right) \quad \text{as } t \rightarrow 0, \quad (7)$$

where  $\alpha > 1$  is not integer and  $n$  is the largest integer smaller than  $\alpha$ .

Since  $\phi(s) = f(t)$ , by Lemma 3 we have  $f_n(t) \sim \phi_n(s)$ , where  $t(s) = (1 - e^{-s}) \sim s$ , as  $s \rightarrow 0$ . So, we can obtain from (7) by using Lemma 1 that

$$\phi_n(s) \sim (-1)^\alpha \Gamma(1 - \alpha) s^\alpha L\left(\frac{1}{s}\right).$$

Now we again apply Theorem 1 to conclude that

$$1 - F_{N(X)} \sim x^{-\alpha} L(x) \text{ as } x \rightarrow \infty.$$

(ii)  $\rightarrow$  (i) Similar to the first part of the proof.  $\square$

Thus, our model for the number of incoming links properly describes an In-Degree distribution that follows a power law with finite expectation and a non-integer exponent.

### 3.3 The main stochastic equation

Combining the ideas from Sections 3.1 and 3.2, we arrive at the following equation

$$R \stackrel{d}{=} c \sum_{j=1}^{N(X)} \frac{1}{d} R_j + (1 - c), \quad (8)$$

where  $c \in (0, 1)$  is the damping factor,  $d \in \{1, 2, \dots\}$  is the fixed Out-Degree of each page, and  $N(X)$  describes the In-Degree of a randomly chosen page as the number of Poisson arrivals on a regularly varying time interval  $X$ . As we discussed above, stochastic equation (8) adequately captures several important aspects of the PageRank distribution and its relation to the In-Degree distribution. Moreover, our model is completely formalized, and thus we can apply analytical methods in order to derive the tail behavior of the random variable  $R$  representing the PageRank.

Linear stochastic equations like (8) have a long history. In particular, (8) is similar to the famous equation that arises in the theory of branching processes and describes many real-life phenomena, for instance, the distribution of the busy period in the  $M/G/1$  queue:

$$B \stackrel{d}{=} \sum_{i=1}^{N(S_1)} B_i + S_1,$$

where  $B$  is the distribution of the busy period (the time interval during which the queue is non-empty),  $S_1$  is the service time of the customer that initiated the busy period,  $N(S_1)$  is the number of Poisson arrivals during this service time and the  $B_i$ 's are independent and distributed as  $B$ . We refer to [23] and other books on queueing theory for more details. Also, see Zwart [24] for an excellent detailed treatment of queues with regular variation, and specifically the busy period problem. We note also that our equation (8) is a special case



in a rich class of stochastic recursive equations that were discussed in detail in the recent survey by Aldous and Bandyopadhyay [3].

This concludes the model description. The next step will be to use our model for providing a rigorous explanation of the indicated connection between the distributions of In-Degree and PageRank.

## 4 Analysis

The idea of our analysis is to write the equation for the Laplace-Stieltjes Transforms of  $X$  and  $R$  and then make use of the Tauberian theorems to prove that  $R$  is regularly varying with the same index as  $X$ . According to Theorem 2, this will give us the desired similarity in tail behavior of the PageRank  $R$  and the In-Degree  $N(X)$ .

As a result of the assumptions from Section 3, we can express the Laplace-Stieltjes transform  $r(s)$  of the PageRank distribution  $R$  in terms of the probability generating function of  $N(X)$  using (8):

$$\begin{aligned} r(s) &:= \mathbb{E}e^{-sR} = e^{-s(1-c)} \mathbb{E} \exp \left( -s \frac{c}{d} \sum_{i=1}^{N(X)} R_i \right) \\ &= e^{-s(1-c)} \sum_{k=1}^{\infty} \mathbb{E} \exp \left( -s \frac{c}{d} \sum_{i=1}^k R_i \right) \mathbb{P}(N(X) = k) \\ &= e^{-s(1-c)} \sum_{k=1}^{\infty} \left( r \left( s \frac{c}{d} \right) \right)^k \mathbb{P}(N(X) = k) \\ &= e^{-s(1-c)} \mathbb{G}_{N(X)} \left( r \left( s \frac{c}{d} \right) \right). \end{aligned}$$

Since, by (6),  $\mathbb{G}_{N(X)}(s) = f(1-s)$ , we arrive at

$$r(s) = f \left( 1 - r \left( \frac{c}{d} s \right) \right) e^{-s(1-c)}. \quad (9)$$

It can be shown (e.g. arguing as in [14, Section XIII.4]) that equation (9) has a unique solution  $r(s)$  which is completely monotone and has  $r(0) = 1$  if and only if  $c/d < 1$ . This inequality is satisfied for the typical values  $d > 1$  and  $0 < c < 1$ .

As in Section 3.2, we will start the analysis with providing the correspondence between existence of the  $n$ -th moments of  $X$  and  $R$ . We remind that  $\xi_1, \dots, \xi_n$  denote the first  $n$  moments of  $X$ . Further, denote the first  $n$  moments of  $R$  by  $\rho_1, \dots, \rho_n$ , and define

$$r_n(s) = (-1)^{n+1} \left( r(s) - \sum_{k=0}^n \frac{\rho_k}{k!} (-s^k) \right),$$

as in (2). Note that taking expectations on both sides of (8) we easily obtain  $\mathbb{E}R = \rho_1 = 1$ . This follows from the independence of  $N(X)$  and the  $R_j$ 's and the fact that  $\mathbb{E}N(X) = \mathbb{E}X = \xi_1 = d$ .

The next lemma holds (the proof is provided in the Appendix).

**Lemma 4.** *The following are equivalent*

- (i)  $\xi_n < \infty$ ,
- (ii)  $\rho_n < \infty$ .

**Remark 3.** *Similar as in Remark 1, we can reformulate Lemma 4 as*

$$f_n(s) = o(s^n) \quad \text{if and only if} \quad r_n(s) = o(s^n).$$

**Remark 4.** *Note that the stochastic inequality  $R \stackrel{d}{>} (1-c) \left( \frac{c}{d} N(X) + 1 \right)$  implies that the tail of the PageRank is at least as heavy as the tail of the In-Degree.*

To establish the main result we only need to make one technical observation, which is proved in the Appendix.

**Corollary 1.** *The following holds:*

$$r_n(s) - dr_n\left(\frac{c}{d}s\right) = f_n(t) + O(t^{n+1}).$$

Now we are ready to explain the similarity between In-Degree and PageRank distributions. The next theorem formalizes this main statement.

**Theorem 3.** *The following are equivalent*

- (i)  $1 - F_{N(X)}(x) \sim x^{-\alpha} L(x) \quad \text{as } x \rightarrow \infty$ ,
- (ii)  $1 - F_R(x) \sim \frac{c^\alpha}{d^\alpha - c^\alpha d} x^{-\alpha} L(x) \quad \text{as } x \rightarrow \infty$ .

*Proof.*

(i)  $\rightarrow$  (ii) From (i) and Theorem 2 it follows that

$$1 - F_X(x) \sim x^{-\alpha} L(x) \quad \text{as } x \rightarrow \infty. \tag{10}$$

Theorem 1 also implies that (10) is equivalent to  $f_n(t) \sim (-1)^\alpha \Gamma(1-\alpha) t^\alpha L\left(\frac{1}{t}\right)$ , where  $t(s) \sim (c/d)s$ , as  $s \rightarrow 0$ . Then, by Corollary 1 we obtain

$$r_n(s) - dr_n\left(\frac{c}{d}s\right) \sim (-1)^n \Gamma(1-\alpha) \left(\frac{c}{d}\right)^\alpha s^\alpha L\left(\frac{1}{s}\right) \quad \text{as } s \rightarrow 0.$$

Then also for every  $k \geq 0$ , as  $s \rightarrow 0$ , we have

$$\begin{aligned} r_n\left(\left(\frac{c}{d}\right)^k s\right) - dr_n\left(\left(\frac{c}{d}\right)^{k+1} s\right) &\sim (-1)^n \Gamma(1-\alpha) \left(\frac{c}{d}\right)^\alpha \left(\frac{c}{d}\right)^{\alpha k} s^\alpha L\left(\frac{1}{\left(\frac{c}{d}\right)^k s}\right) \\ &\sim (-1)^n \Gamma(1-\alpha) \left(\frac{c}{d}\right)^\alpha \left(\frac{c}{d}\right)^{\alpha k} s^\alpha L\left(\frac{1}{s}\right). \end{aligned}$$

Using the infinite-sum representation for  $r_n(s)$ ,

$$r_n(s) = \sum_{k=0}^{\infty} d^k \left( r_n \left( \left( \frac{c}{d} \right)^k s \right) - d r_n \left( \left( \frac{c}{d} \right)^{k+1} s \right) \right),$$

(see the proof of Lemma 4 in the Appendix, equation (17)), we directly obtain

$$r_n(s) \sim (-1)^n \Gamma(1 - \alpha) \frac{d^\alpha}{d^\alpha - c^\alpha d} \left( \frac{c}{d} \right)^\alpha s^\alpha L \left( \frac{1}{s} \right) \quad \text{as } s \rightarrow 0.$$

Now we again apply Theorem 1, which leads to (ii).

(ii)  $\rightarrow$  (i) The proof follows easily from (ii) and Corollary 1.  $\square$

Thus, we have shown that the asymptotic behaviors of PageRank and In-Degree differ only by the multiplicative factor  $\frac{c^\alpha}{d^\alpha - c^\alpha d}$ , while the power law exponent remains the same. In the next section we will experimentally verify this result.

Note that in the present model, we can easily account for the pages without out-going links (*dangling nodes*). There are different ways to deal with such nodes when defining the ranking. We consider a classical scenario, where a dangling node is equivalent to a node, that has an outgoing link to every page in the Web. Then equation (1) becomes

$$PR(i) = c \sum_{j \rightarrow i} \frac{1}{d_j} PR(j) + \frac{c}{n} \sum_{j \in \mathcal{D}} PR(j) + (1 - c), \quad (11)$$

where  $\mathcal{D}$  is a set of dangling nodes, and  $n$  is the number of pages in the Web. Assume further that the PageRank of a random page does not depend on the fact whether the page is dangling. Indeed, it can be shown that the PageRank of a page can not be altered significantly by modifying outgoing links [5]. Moreover, experiments e.g. in [13] show that dangling nodes are often just regular pages whose links have not been crawled, for instance, because it was not allowed by `robot.txt`. Besides, even authentically dangling pages such as `.pdf` or `.ps` files, often contain important information and gain high ranking independently of the fact that they do not have outgoing links. Such independence implies, in particular, that the average PageRank of dangling nodes is 1, and thus the fraction of the total PageRank mass concentrated in dangling nodes, equals to the fraction of dangling nodes  $p_0$ :

$$p_0 = \frac{|\mathcal{D}|}{n} = \frac{1}{n} \sum_{j \in \mathcal{D}} PR(j).$$

Note that in the presence of dangling nodes, the average out-degree of non-dangling nodes is  $d(1 - p_0)^{-1}$ .

Now, exactly as our main stochastic equation (3) is analogous to (1), we can also provide a stochastic equation analogous to (11) as follows:

$$R \stackrel{d}{=} c \sum_{j=1}^N \frac{1 - p_0}{d} R_j + [1 - c(1 - p_0)].$$

Observe that this is the same stochastic equation as (3), only with  $c(1 - p_0)$  instead of  $c$ . Thus, Theorem 3 applies directly after the corresponding straightforward adjustments. Since our data set contains a negligible amount of dangling nodes, we do not take them into account in the experiments.

## 5 Numerical Results

### 5.1 Power Law Identification

The identification and measuring of power law behavior is not always simple. In this section we provide a brief overview of techniques that we used to plot and numerically identify power law distributions.

The standard strategy is to plot a histogram of a quantity on logarithmic scales to obtain a straight line, which is a typical feature of the power law. However, this technique is often not efficient. In [21], Newman clearly illustrated that even for generated random numbers with a known distribution the noise in the tail region has a strong influence on the estimation of the power law parameters. He suggests to plot the fraction of measurements that are not smaller than a given value, i.e. the complementary cumulative distribution function  $1 - F(x) = \mathbb{P}(X > x)$  rather than the histogram. The advantage is that we obtain a less noisy plot. Besides, this idea is consistent with our analysis in the previous section, which was based on complementary cumulative distribution functions. We note that if the distribution of  $X$  follows a power law with exponent  $\alpha$  so that  $1 - F(x) \sim Cx^{-\alpha}$ ,  $x \rightarrow \infty$ , where  $C$  is some constant, then the corresponding histogram has an exponent  $\alpha + 1$ . Thus, the plot of  $1 - F(x)$  on logarithmic scales has a smaller slope than the plot of the histogram.

Computing the correct slope from the observed data is also not trivial. Goldstein et al. in [17], and later Newman in [21], have proposed to use a maximum likelihood estimator, which provides a more robust estimation of the power law exponent than the standard least-squares fit method. Thus, we compute the exponent  $\alpha$  using the next formula from [21]:

$$\alpha = 1 + N \left( \sum_{i=1}^N \ln \frac{x_i}{x_{min}} \right)^{-1}. \quad (12)$$

Here the quantities  $x_i$ ,  $i = 1, \dots, N$ , are the measured values of  $X$ , and  $x_{min}$  usually corresponds to the smallest value of  $X$  for which the power law behavior is assumed to hold.

In the next sections we will present our experiments on real Web Data and on a graph that represents a well-known mathematical model of the Web (Growing Networks). In both cases, for each value  $x$ , we plot in log-log scale the fraction of measurements that are not smaller than  $x$ , and we use (12) to obtain the exponents.

## 5.2 Web Data

To confirm our results on asymptotic similarity between PageRank and In-Degree distributions we performed experiments on the public data of the Stanford Web from [1]. We calculated all PageRank values for a Web graph with 281903 nodes (pages) and  $\sim 2.3$  million edges (links) using the standard power method (see e.g. [19]). For this data set, the average Out-Degree, and hence average In-Degree is 8.2.

There are several papers, see [6], [16], [12] and [22], that describe similar experiments for different domains and different number of pages, and they all confirm that PageRank and In-Degree follow power laws with the same exponent, around 2.1. In Figure 1 we show the log-log plots for In-Degree and PageRank of the Stanford Web Data, for different values of the damping factor ( $c = 0.1, 0.5$  and  $0.9$ ). Clearly, these empirical values of In-Degree and PageRank constitute parallel straight lines for all values of the damping factor, provided that the PageRank values are reasonably large. It was observed in [6] that in general, PageRank depends on the damping factor but the PageRank of the top 10% of pages obeys a power law with the same exponent as the In-Degree, independent on the damping factor. This is in perfect agreement with our experimental results and the mathematical model, which is focused on the right tail behavior of the PageRank distribution.

The calculations based on the maximum likelihood method yield a slope  $-1.1$ , which verifies that In-Degree and PageRank have power laws with the same exponent  $\alpha = 1.1$  (this corresponds to the well known value 2.1 for the histogram). More precisely, in Figure 1 we fitted the lines  $y = -1.1x + 0.08$ ,  $y = -1.1x - 0.87$ ,  $y = -1.1x - 1.27$ , and  $y = -1.1x - 2.07$  to the plots of In-Degree and PageRank (with  $c = 0.9$ ,  $c = 0.5$  and  $c = 0.1$ , respectively).

We also investigated whether Theorem 3 correctly predicts the multiplicative factor

$$y(c) = \frac{c^\alpha}{d^\alpha - c^\alpha d}.$$

In Figure 2 we plotted  $\log_{10}(y(c))$  and we compared it to the observed differences between the logarithms of the complementary cumulative distribution functions of PageRank and In-Degree, for different values of the damping factor. Again  $d = 8.2$  because that is the average the average In/Out-Degree in the data set. As can be seen, the theoretical and observed values are quite close. E.g., for typical values of  $c$  between 0.8 and 0.9, the difference is 0.41, resulting in a factor  $y(c)$  that is only a factor 2.57 larger than in the observed data. Thus, our model not only allows to prove the similarity in the power law behavior but also gives a good approximation for the difference between the two distributions.

The discrepancy between the predicted and observed values of the multiplicative factor suggests that our model does not capture PageRank behavior to the full extent. For instance, the assumption of the independence of PageRank values of pages that have a common neighbor may be too strong. We believe however that the achieved precision, especially for small values of  $c$ , is quite good for our relatively simple stochastic model.

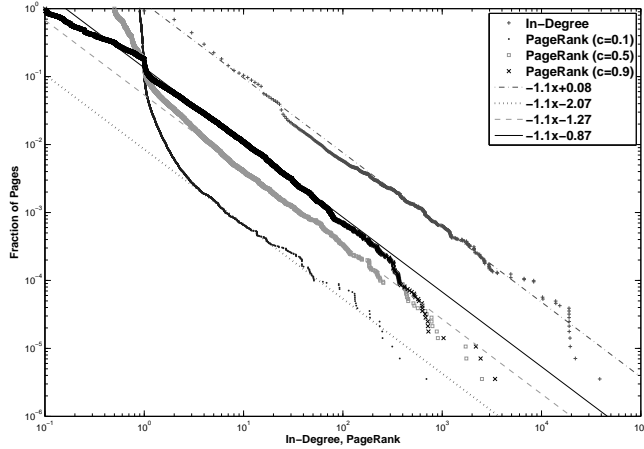


Figure 1: Plots for the Web data. Number of pages with In-Degree/PageRank greater than  $x$  versus  $x$  in log-log scale, and the fitted straight lines.

### 5.3 Growing Networks

Growing Networks, introduced by Barabási and Albert [2], now represent a large class of models that are commonly accepted as a possible scenario of Web growth. In particular, these models provide a mathematical explanation for the power law behavior of In-Degree [10]. The recent studies [4], [15] addressed for the first time the PageRank distribution in Growing Networks.

Growing Network models are characterized by preferential attachment. This entails that a newly created node connects to the existing nodes with probabilities that are proportional to the current In-Degrees of the existing nodes. We simulated a slightly modified version of this model, where a new link points to a randomly chosen page with probability  $\beta$ , and with probability  $1 - \beta$  the preferential attachment selection rule is used. This allows us to tune the exponent of the resulting power law [21].

We simulate our Growing Network using Matlab. We start with  $d$  nodes and at each step we add a new node that links to  $d$  already existing nodes. To ensure the same number of outgoing links for all pages, at the end of the simulation, we link the first  $d$  nodes to randomly chosen pages. In the example presented below we set  $\beta = 0.2$  and obtain a network of 50000 nodes with Out-Degree  $d = 8$ .

In Figure 3 we present the numerical data for In-Degree and PageRank in the Growing Network. Clearly, the Web data from Section 5.2 shows a much better agreement with our model than the data generated by the preferential attachment algorithm. In the next section we briefly compare recent results on PageRank in Growing Networks to our present study and we indicate possible directions for further research.

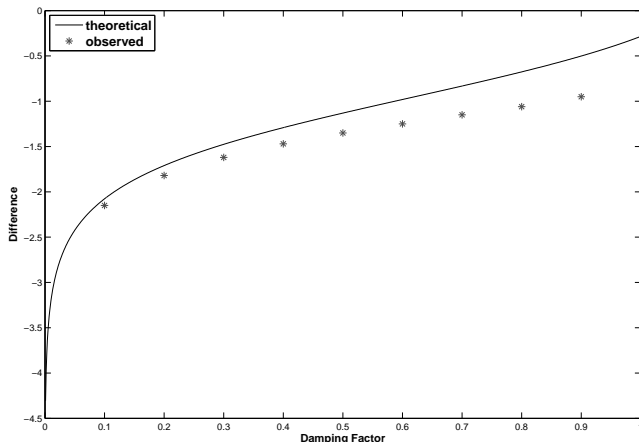


Figure 2: The theoretical and observed differences between logarithmic asymptotics of In-Degree and PageRank.

## 6 Discussion

Our model and analysis resulted in the conclusion that PageRank and In-Degree should follow power laws with the same exponent. Growing Network models may provide an alternative explanation [4, 15]. For instance, in the recent paper by Avrachenkov and Lebedev [4] it was shown that the *expected* PageRank in Growing Networks follows a power law with an exponent, which does depend on the damping factor but equals  $\approx 2.08$  for  $c = 0.85$ . Thus, the model in [4] can also be used to explain the tail behavior of PageRank, but it leads to a slightly different result than our model because in our case the power law exponent of PageRank does *not* depend on the damping factor. The reason could be that we focus only on the asymptotics, whereas [4] employs a mean-field approximation. Indeed, experiments show that the shape of the PageRank distribution does depend on the damping factor, and thus, it may affect the average values, whereas the tail behavior remains the same for all values of  $c$ .

We emphasize that compared to [4, 15], our model provides a completely different approach for modelling the relation between In-Degree and PageRank. Specifically, we do not make any assumption on the underlying Web graph, whereas [4, 15] choose the preferential attachment structure, thus exploiting the fact that this graph model correctly captures the In-Degree distribution. We believe that both approaches should be elaborated and used in further research on the PageRank distribution.

One of the important innovations in the present work is the analogy between the PageRank equation and the equation for the busy period that enables us to apply the techniques from [20]. In fact, queueing systems with heavy tails and in

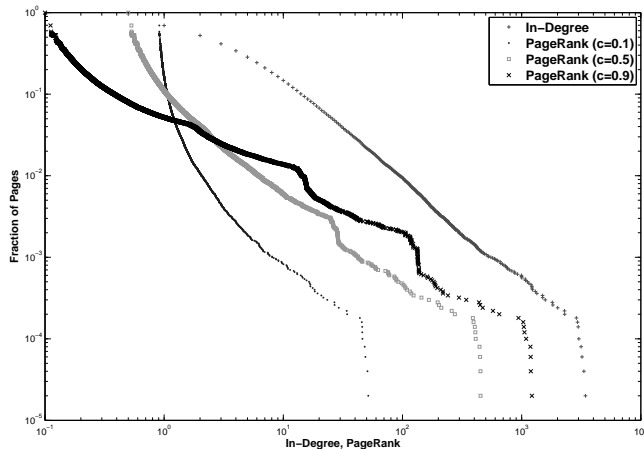


Figure 3: Plots for the Growing Network model. Number of pages with In-Degree/PageRank greater than  $x$  versus  $x$  in log-log scale.

particular the busy period problem allow for a more sophisticated probabilistic analysis (see e.g. [24]). It would be interesting to apply these advanced methods to the problems related to the World Wide Web and PageRank.

Our model definitely lacks the dependencies between the PageRank values of pages sharing a common neighbor. Such dependencies must be present in the Web in particular due to the high clustering of the Web graph [21] (roughly speaking, clustering means that with high probability, two neighbors of the same page are connected to each other). Thus, in our further research we could try to include some sort of dependencies in our stochastic equation. Another natural way to bring our model closer to the real-life situation is to allow random Out-Degrees. Besides, we could also consider personalization or topic sensitivity [18]. The impact of these factors on the PageRank distribution could be determined by extending and generalizing the proposed analytical model.

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## Appendix

*Proof of Lemma 3.*

(i)  $\rightarrow$  (ii) By Lemma 2 we know that  $\xi_n < \infty$  if and only if  $f(t)$  can be written as

$$f(t) = \sum_{i=0}^n \frac{\xi_i}{i!} (-t)^i + o(t^n) \text{ as } t \rightarrow 0.$$

Denote  $t(s) := 1 - e^{-s}$ , then  $t(s) \rightarrow 0$  as  $s \rightarrow 0$ , and we can substitute

$$\begin{aligned} \phi(s) &= f(1 - e^{-s}) = \sum_{i=0}^n \frac{\xi_i}{i!} (-(1 - e^{-s}))^i + o((1 - e^{-s})^n) \\ &= \sum_{i=0}^n \frac{\xi_i}{i!} (-1)^i \left( \sum_{k=1}^{\infty} (-1)^{k+1} \frac{s^k}{k!} \right)^i + o(s^n), \end{aligned}$$

which can be written as

$$\phi(s) = \sum_{i=0}^n \frac{\nu_i}{i!} (-s)^i + o(s^n)$$

for some finite constants  $\nu_0 = 1$  and  $\nu_1, \dots, \nu_n$ , that can be expressed in terms of  $\xi_0 = 1$  and  $\xi_1, \dots, \xi_n$ . Thus, by uniqueness of the power series expansion and by Lemma 2 we have  $\nu_n < \infty$ .

(ii)  $\rightarrow$  (i) Similarly,  $s(t) := -\ln(1-t) \rightarrow 0$  as  $t \rightarrow 0$ , so we obtain

$$\begin{aligned} f(t) &= \phi(-\ln(1-t)) = \sum_{i=0}^n \frac{\nu_i}{i!} \ln^i(1-t) + o(\ln^n(1-t)) \\ &= \sum_{i=0}^n \frac{\nu_i}{i!} \left( -\sum_{k=1}^{\infty} \frac{t^k}{k} \right)^i + o\left( \left( -\sum_{k=1}^{\infty} \frac{t^k}{k} \right)^n \right) \\ &= \sum_{i=0}^n \frac{\xi_i}{i!} (-t)^i + o(t^n), \end{aligned}$$

for  $\xi_0 = 1$  and some  $\xi_1, \dots, \xi_n$  that can be expressed in terms of  $\nu_0 = 1$  and  $\nu_1, \dots, \nu_n$ , which similarly implies  $\xi_n < \infty$ .  $\square$

*Proof of Lemma 4.*

(i)  $\rightarrow$  (ii) We use induction, starting from  $n = 1$  for which both (i) and (ii) are valid. Assume that for  $k = 1, 2, \dots, n-1$  it has been shown that (i)  $\rightarrow$  (ii). We introduce the following notation, to be used throughout this section. Denote

$$\begin{aligned} g(s) &:= e^{-s(1-c)}, \quad \text{and} \\ t(s) &:= 1 - r\left(\frac{c}{d}s\right). \end{aligned}$$

Then we can write (9) as

$$r(s) = f(t)g(s). \quad (13)$$

We know from (i) that

$$\begin{aligned} f(t) &= 1 - dt + \sum_{k=2}^n \frac{\xi_k (-t)^k}{k!} + o(t^n) \\ &= 1 - d \left( 1 - r\left(\frac{c}{d}s\right) \right) + \sum_{k=2}^n \frac{\xi_k (-t)^k}{k!} + o(t^n). \end{aligned}$$

Thus, from (13) we obtain

$$r(s) - dg(s)r\left(\frac{c}{d}s\right) = \left( 1 - d + \sum_{k=2}^n \frac{\xi_k (-t)^k}{k!} + o(t^n) \right) g(s). \quad (14)$$

However, it follows from the induction hypothesis for  $n-1$  that

$$r(s) = 1 - s + \sum_{k=2}^{n-1} \frac{\rho_k}{k!} (-s^k) + o(s^{n-1}),$$

so we can present  $t(s)$  as a sum

$$t(s) = - \sum_{k=1}^{n-1} \frac{\rho_k}{k!} \left( \frac{c}{d} \right)^k (-s)^k + o(s^{n-1}).$$

Using this, we can actually find  $t^k(s)$ :

$$t^k(s) = \sum_{i=k}^{n+k-2} \beta_{k,i} s^i + o(s^{n+k-2}), \quad (15)$$

for  $k \geq 1$  and appropriate constants  $\beta_{k,i}, i = k, \dots, k+n-2$ . Thus, we obtain by (14) and (15):

$$r(s) - dg(s)r\left(\frac{c}{d}s\right) = \left(\sum_{i=0}^n \gamma_i (-s)^i + o(s^n)\right) g(s)$$

for appropriate constants  $\gamma_0, \dots, \gamma_n$ . Using the expansion of  $g(s)$ , it is not difficult to show that for appropriate constants  $\eta_0, \dots, \eta_n$ , we also have

$$r(s) - dr\left(\frac{c}{d}s\right) = \sum_{i=0}^n \eta_i s^i + o(s^n).$$

In other words, because of the uniqueness of the series expansion, we have

$$\left(r(s) - dr\left(\frac{c}{d}s\right)\right)_n = r_n(s) - dr_n\left(\frac{c}{d}s\right) = o(s^n). \quad (16)$$

We will now show that this implies (ii), to which end we consider the partial sums

$$\begin{aligned} r_n^N(s) &= \sum_{k=0}^N d^k \left( r_n \left( \left( \frac{c}{d} \right)^k s \right) - dr_n \left( \left( \frac{c}{d} \right)^{k+1} s \right) \right) \\ &= r_n(s) - d^{N+1} r_n \left( \left( \frac{c}{d} \right)^{N+1} s \right). \end{aligned}$$

Taking the limit as  $N \rightarrow \infty$ , we have for the last term that

$$\begin{aligned} &\lim_{N \rightarrow \infty} d^{N+1} r_n \left( \left( \frac{c}{d} \right)^{N+1} s \right) \\ &= \lim_{N \rightarrow \infty} \frac{r_n \left( \left( \frac{c}{d} \right)^{N+1} s \right)}{\left( \left( \frac{c}{d} \right)^{N+1} s \right)^{n-1}} \lim_{N \rightarrow \infty} \left( \frac{c}{d} \right)^{(N+1)(n-2)} s^{n-1} c^{N+1} = 0, \end{aligned}$$

where we used the induction hypothesis  $r_n(s) = o(s^{n-1})$  together with  $n \geq 2$ ,  $0 < c < 1$  and  $d > 1$ . It follows that we can express  $r_n(s)$  as an infinite sum,

$$r_n(s) = \sum_{k=0}^{\infty} d^k \left( r_n \left( \left( \frac{c}{d} \right)^k s \right) - dr_n \left( \left( \frac{c}{d} \right)^{k+1} s \right) \right), \quad (17)$$

where we can apply (16) to each of the terms. Further, by definition of  $o(s^n)$ , for all  $\varepsilon > 0$ , there exists a  $\delta = \delta(\varepsilon)$  such that  $|r_n(s) - dr_n(\frac{c}{d}s)| < \varepsilon s^n$  whenever

$0 < s \leq \delta$ . Moreover, for this  $\varepsilon$  and  $\delta$ , and  $0 < s \leq \delta$ , we also have

$$\begin{aligned}
|r_n(s)| &= \left| \sum_{k=0}^{\infty} d^k \left( r_n \left( \left( \frac{c}{d} \right)^k s \right) - dr_n \left( \left( \frac{c}{d} \right)^{k+1} s \right) \right) \right| \\
&\leq \sum_{k=0}^{\infty} \left| d^k \left( r_n \left( \left( \frac{c}{d} \right)^k s \right) - dr_n \left( \left( \frac{c}{d} \right)^{k+1} s \right) \right) \right| \\
&< \sum_{k=0}^{\infty} \varepsilon d^k \left( \frac{c}{d} \right)^{kn} s^n = \frac{d^{n-1} - c^n}{d^{n-1} - c^n} \varepsilon s^n.
\end{aligned} \tag{18}$$

Here the second inequality holds because  $0 < \left( \frac{c}{d} \right)^k s \leq \delta$  for every  $k \geq 0$ . Since for every  $\varepsilon_0 > 0$  there exists  $\delta_0$  such that

$$\left| r_n(s) - dr_n \left( \frac{c}{d} s \right) \right| < \frac{d^{n-1} - c^n}{d^{n-1} - c^n} \varepsilon_0 s^n$$

for  $0 < s \leq \delta_0$ , then according to (18), we have  $|r_n(s)| < \varepsilon_0 s^n$  whenever  $0 < |s| \leq \delta_0$ , by which we have shown that  $r_n = o(s^n)$ .

(ii)  $\rightarrow$  (i) Assume that there exists a nonnegative random variable  $R$  satisfying (8). Then, obviously,  $R \geq 1 - c$ . Moreover, (8) also implies that  $R$  is stochastically greater than  $(1 - c) \left( \frac{c}{d} N(X) + 1 \right)$ . Hence, the existence of the  $n$ -th moment of  $R$  ensures the existence of the  $n$ -th moment of  $N(X)$ , which in turn by Lemma 3 ensures the existence of the  $n$ -th moment of  $X$ .  $\square$

*Proof of Corollary 1.* The proof follows from the first part of the proof of Lemma 4. By definitions of  $r_n(s)$ ,  $f_n(t)$ ,  $t(s)$  and Lemma 4, it follows from (13) that for fixed  $n$ ,

$$\begin{aligned}
(-1)^{n+1} r_n(s) + \sum_{k=0}^n \frac{\rho_k}{k!} (-s^k) &= \left( (-1)^{n+1} f_n(t) + 1 - dt + \sum_{k=2}^n \frac{\xi_k (-t)^k}{k!} \right) g(s) \\
&= \left( (-1)^{n+1} f_n(t) + 1 - d + d \left( (-1)^{n+1} r_n \left( \frac{c}{d} s \right) + \sum_{k=0}^n \frac{\rho_k}{k!} \left( \frac{c}{d} \right)^k (-s)^k \right) + \right. \\
&\quad \left. + \sum_{k=2}^n \frac{\xi_k (-t)^k}{k!} \right) (1 + o(1)).
\end{aligned}$$

Because  $r_n(s) = o(s^n)$  we can extend (15) for  $k \geq 1$  and appropriate constants  $\beta_{k,i}$ ,  $i = k, \dots, k+n-1$ :

$$t^k(s) = \sum_{i=k}^{n+k-1} \beta_{k,i} s^i + o(s^{n+k-1}),$$

and rewrite the last equation as

$$(-1)^{n+1} r_n(s) + \sum_{k=0}^n \frac{\rho_k}{k!} (-s^k)$$

$$= (-1)^{n+1} f_n(t) - d(-1)^{n+1} r_n \left( \frac{c}{d} s \right) + \sum_{k=0}^{n+1} \tau_k s^k + o(s^{n+1}),$$

where  $\tau_0, \dots, \tau_{n+1}$  are corresponding constants. Now due to the uniqueness of the series expansion, we can reduce the above formula to

$$r_n(s) = f_n(t) + dr_n \left( \frac{c}{d} s \right) + (-1)^{n+1} \tau_{n+1} s^{n+1} + o(s^{n+1}).$$

Then we get:

$$r_n(s) - dr_n \left( \frac{c}{d} s \right) = f_n(t) + O(s^{n+1}).$$

□